

All of us after next week

PSO 14

K-D Trees, Point Trees



Announcements

1. Fill out the instructor feedback surveys (ty 40% of you)
2. Last review session on Friday (time TBD, location TBD (existence TBD))
3. No OH next week
 - a. I will NOT be on duty

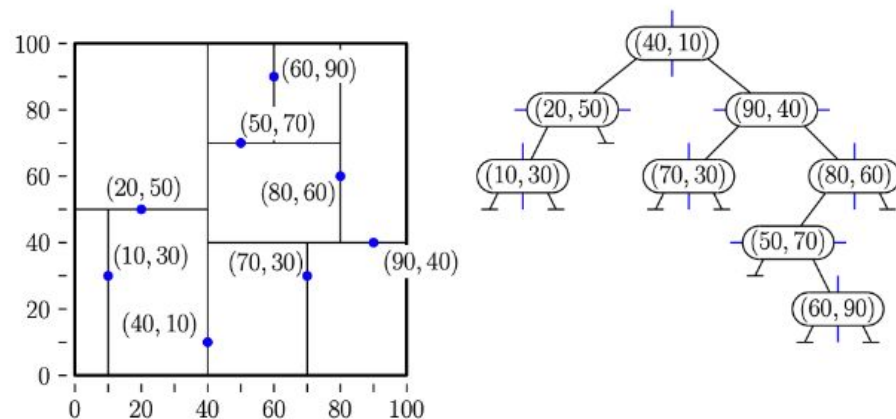
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2. Last review session on Friday (time TBD, location TBD (existence TBD))
3. No OH next week
 - a. I will NOT be on duty but..
 - b. I *might* happen to be sitting around the commons from 12-2PM Sat,Sun,Mon,Tues
 - c. I *might* be open to answering any questions if they *happen* to be asked
 - d. I *might* be hungover

Question 1

(kd trees)

(1) Consider the kd-tree shown in the figure below. Assume a standard kd-tree where the cutting dimensions alternates between x and y with each level.



- (1) Show the final tree after the operation **insert**((70,50)).
- (2) Starting with the original tree, show the final tree after **delete**((40,10)).
- (3) Starting with the original tree, show the final tree after **delete**((80,60)).

Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

- (a) What is the expected **depth** $D(n)$ of the QuadTree?
- (b) What is the total number of **leaf nodes** in the tree in terms of n ?
- (c) If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

But first..

Question 3

(Backward pattern matching)

The Boyer-Moore algorithm is based upon backward pattern matching. Let us do a simple review via the following questions:

1. Run Boyer-Moore algorithm in the following worst-case scenario:

$$T := \underbrace{aaa \cdots a}_9 \quad \text{and} \quad P := baaaaa.$$

Boyer-Moore: Iteratively compare pattern P with target, going backward

T	a	a	a	a	a	a	a	a
P	b	a	a	a	a	a		

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T[0] does not equal P[0]! Next steps..

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T[0] does not equal P[0]! Next steps.. We mismatched on target **a**
The last occurrence of pattern **a**

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P	b	a	a	a	a	a		

Move P (to align target **a** with pattern **a**) OR (one after target mismatch)

Whichever moves P the *least* amount – in this ex. We move one after mismatch

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Fast forward..

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Fast forward.. Same mismatch, jump 1

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Boyer-Moore: Iteratively compare pattern P with target, going backward

T	a	a	a	a	a	a	a	a
P			b	a	a	a	a	

Same thing will happen 1 more time

The example from last time

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

The example from last time

Mismatch here

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

The example from last time

We mismatched on target **X**

The last occurrence of pattern **X**

T	O	O	X	X	X	X	O	O	O
P	O	X	X	X	X	O	O	O	

Move P (to align target **X** with pattern **X**) OR (one after target mismatch)
Whichever moves P the *least* amount

The example from last time

We mismatched on target **X**

The last occurrence of pattern **X**

T	O	O	X	X	X	X	O	O	O
P		O	X	X	X	X	O	O	O

Move P (to align target **X** with pattern **X**) OR (one after target mismatch)

Whichever moves P the *least* amount

Last note: If there is no last occurrence of the target mismatch, default to one after mismatch

View this at your leisure – a longer example

<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>	<i>x</i>	<i>x</i>	<i>b</i>										
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>				<i>x</i>	<i>x</i>	<i>b</i>							
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>					<i>x</i>	<i>x</i>	<i>b</i>						
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>								<i>x</i>	<i>x</i>	<i>b</i>			
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>									<i>x</i>	<i>x</i>	<i>b</i>		
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>										<i>x</i>	<i>x</i>	<i>b</i>	
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
<i>P</i>										<i>x</i>	<i>x</i>	<i>b</i>	

The green is a comparison made

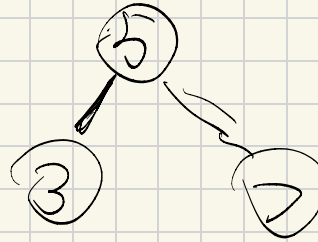
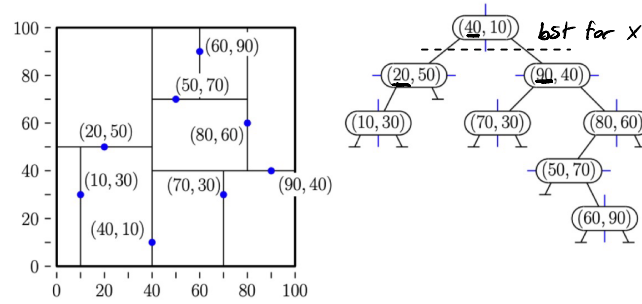
View this at your leisure – a longer example

c not in the pattern
We move one after
(In this case, big jump)

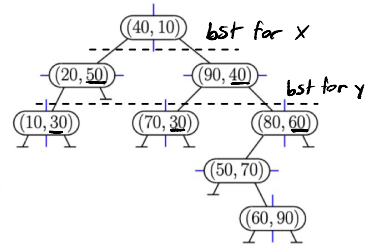
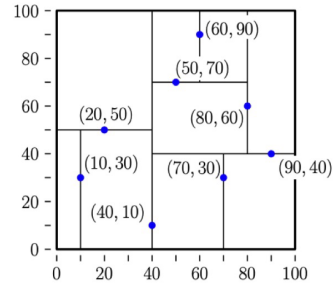
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
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<i>P</i>					<i>x</i>	<i>x</i>	<i>b</i>						
<i>T</i>	<i>a</i>	<i>b</i>	<i>c</i>	<i>a</i>	<i>b</i>	<i>x</i>	<i>a</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>x</i>	<i>b</i>
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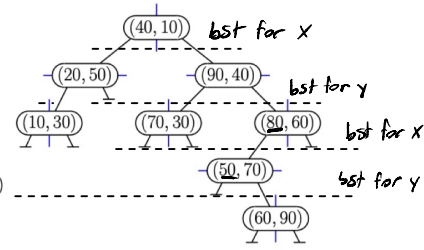
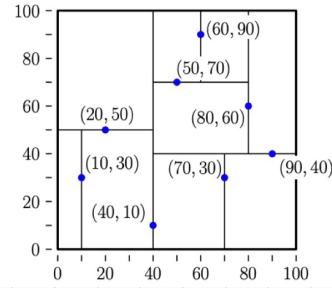
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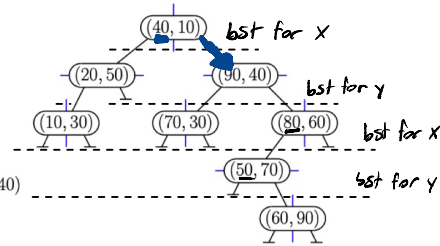
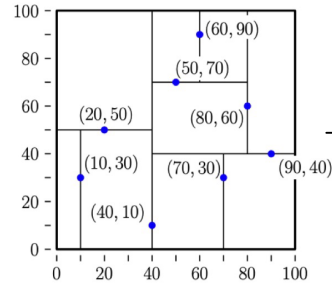


(2) **(kd-trees)** Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation $\text{insert}((70,50))$.



Now let's insert (70,50)

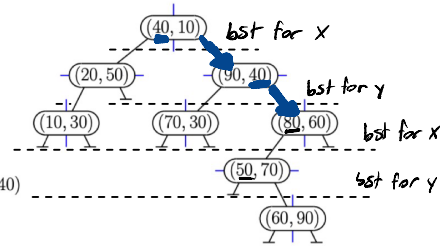
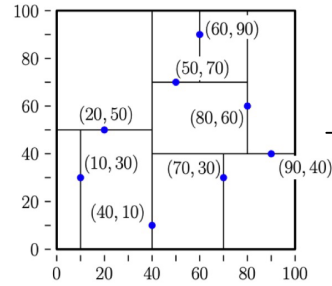
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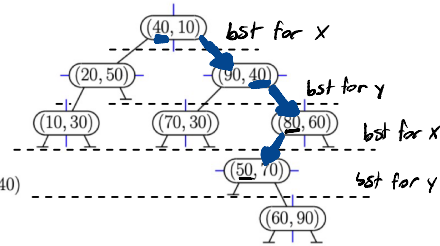
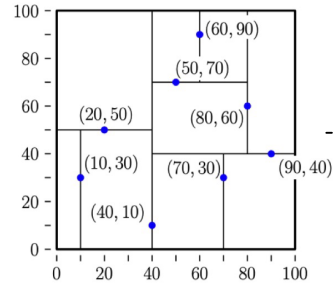
6

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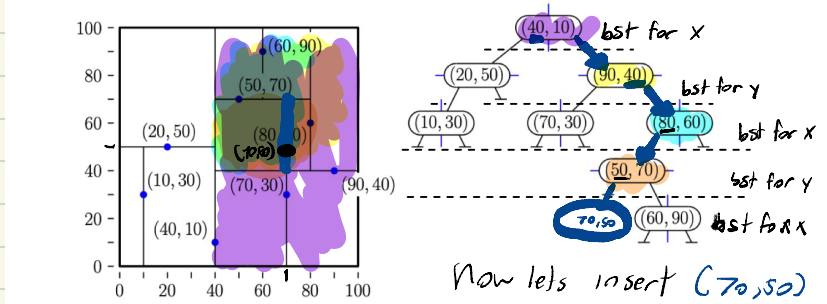
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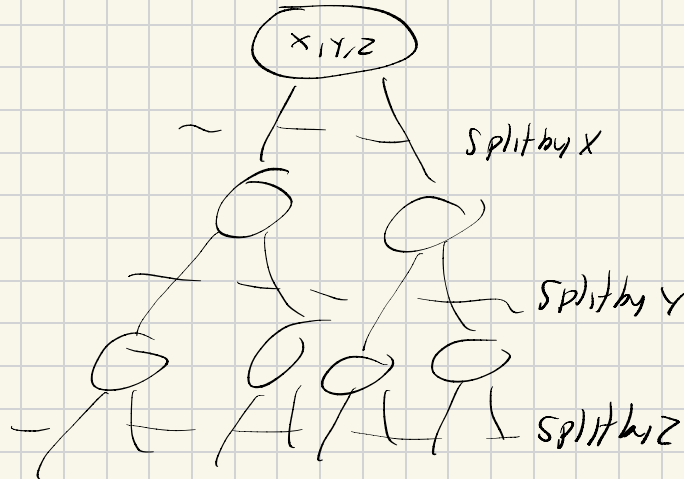
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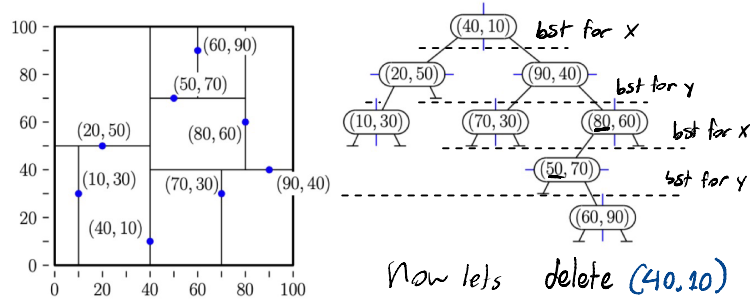


2d tree

3d tree

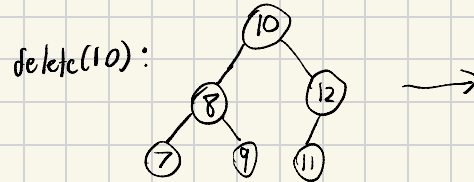


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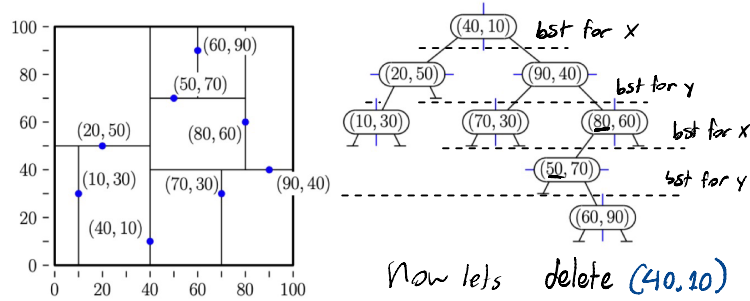


Recall deletion in normal BST

- Find predecessor/successor leaf to replace
ex:

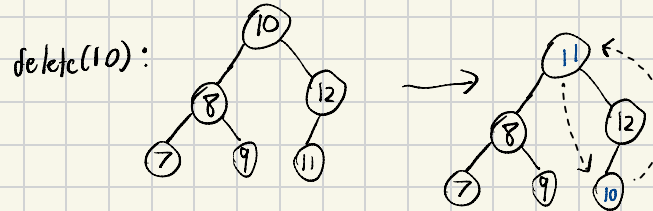


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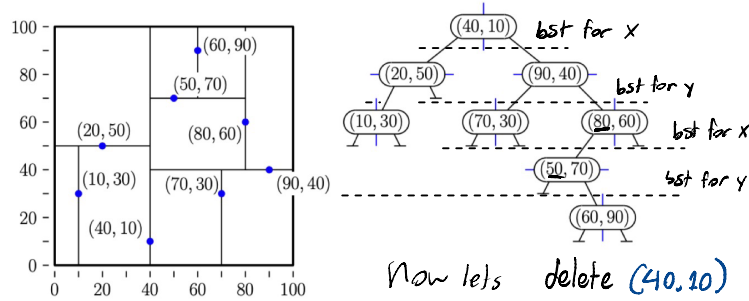


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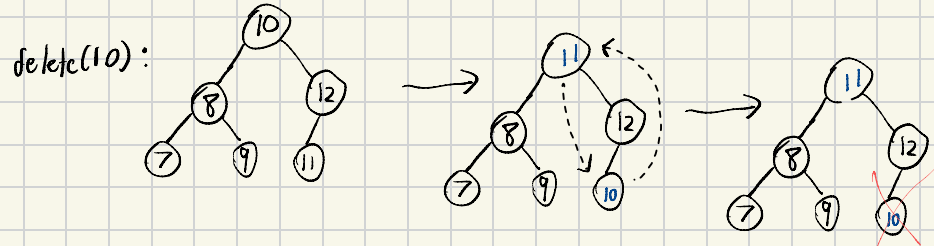


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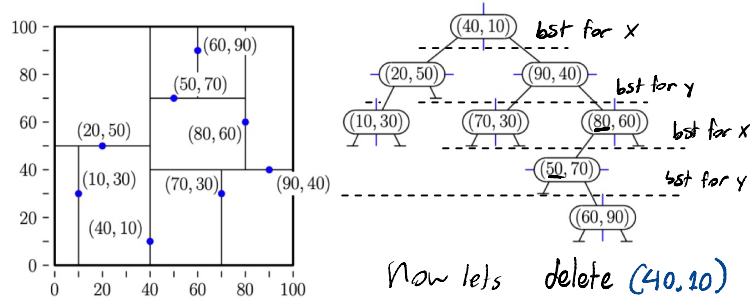
Recall deletion in normal BST

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ex:



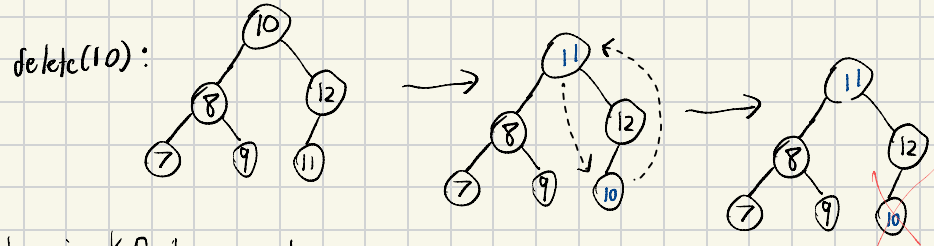
(If not a + leaf continue recursively)

(2) (**kd-trees**) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation $\text{insert}((70,50))$.



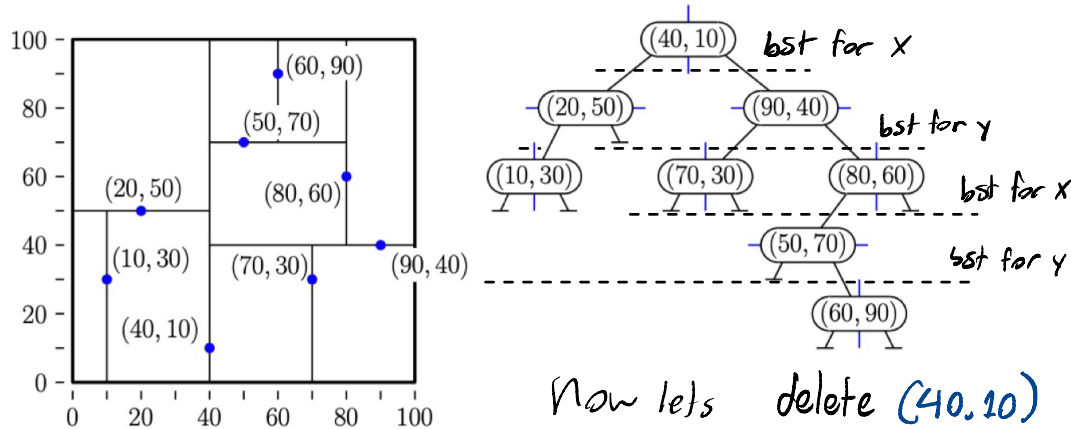
Recall deletion in normal BST

- Find predecessor/successor leaf to replace
ex:



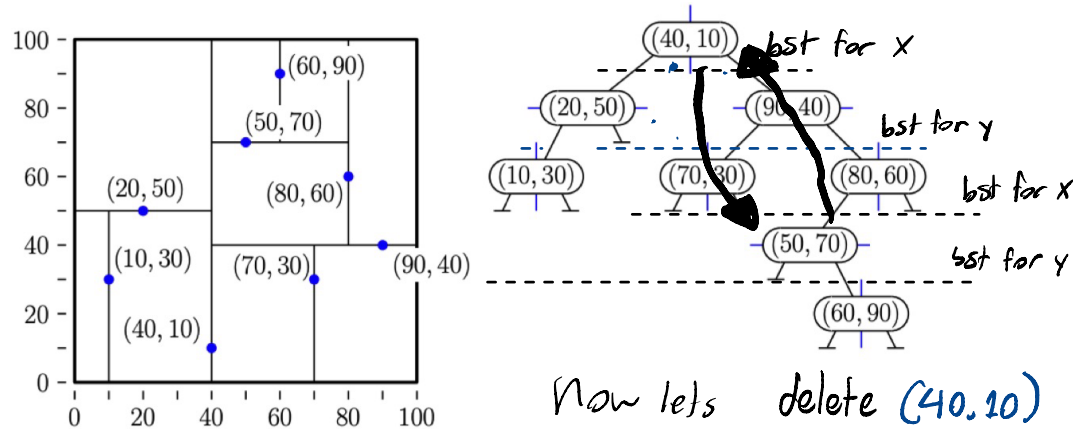
Deletion in KD-trees is the same,
except you choose successor/predecessor based on *dimension*

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1) Find successor for $X=40$

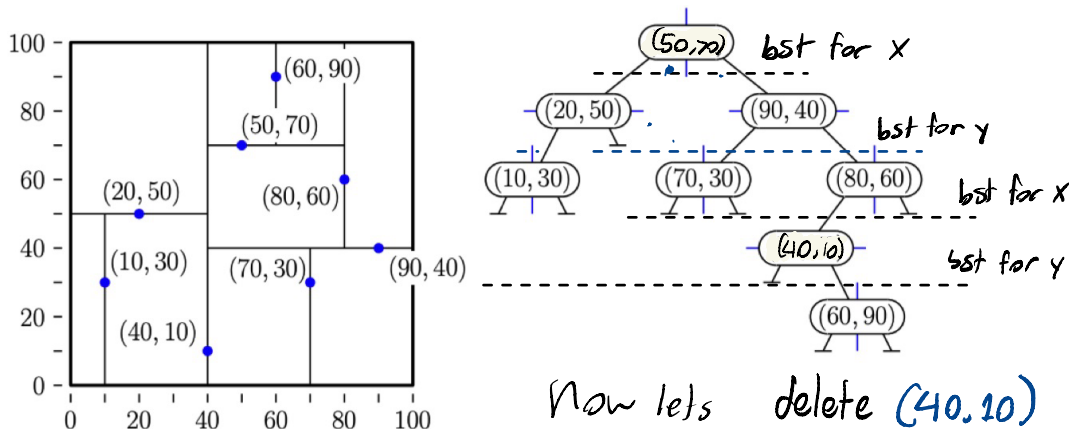
(2) (**kd-trees**) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation `insert((70,50))`.



2) Swap with (40,10)

[Think about why this preserves k-d order]

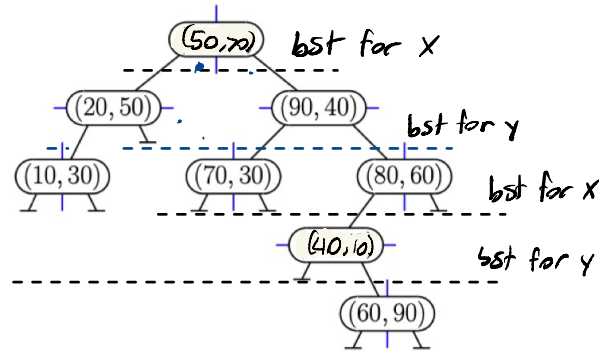
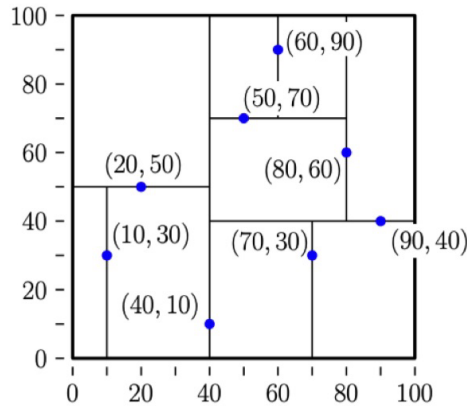
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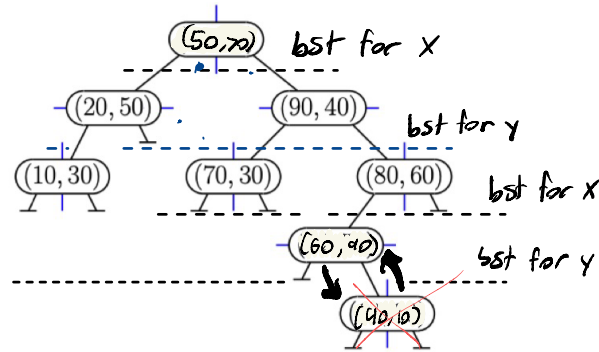
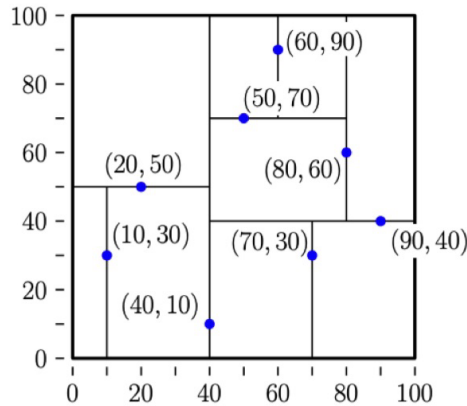
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Now lets delete (40,10)

2.2) Not a leaf yet, find successor in $y=10$ and swap.

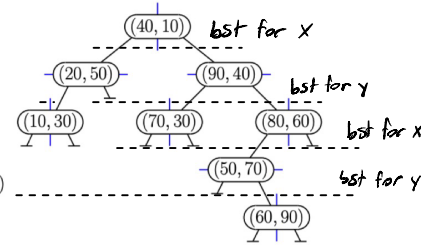
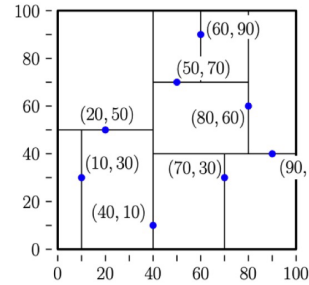
(2) (**kd-trees**) Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation $\text{insert}((70,50))$.



Now lets delete (40,10)

2.2) Not a leaf yet, find successor in $y=10$ and swap.

(2) **(kd-trees)** Consider the kd-tree shown below. We assume it's a standard kd-tree where the cutting dimensions alternates between x and y with each level. Show the final tree after the operation $\text{insert}((70,50))$.



Now let's delete (80,60).

Left as
Exercise

NOTE: If no successor, find a predecessor instead and vice versa.

Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

- (a) What is the expected **depth** $D(n)$ of the QuadTree?
- (b) What is the total number of **leaf nodes** in the tree in terms of n ?
- (c) If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

for now

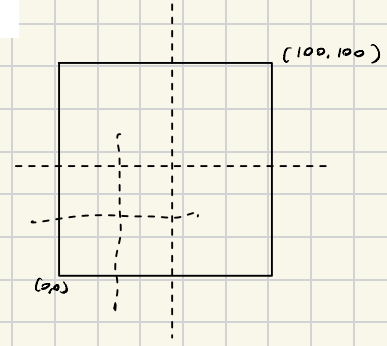
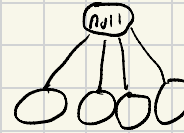
I'm going to ignore this question[^] and go over how we insert into a quad tree.

Question 1

(1) **(Quadrees)** Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

• Suppose $root.pt = null$
 $root.region = (0, 0, 100, 100)$
 $root.children = []$

1) Insert (35,40)



Question 1

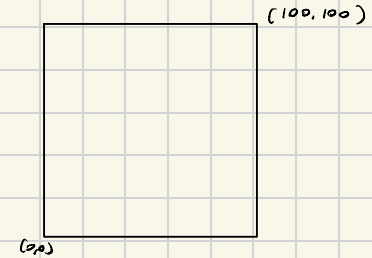
(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

• Suppose $root.pt = null$
 $root.region = (0, 0, 100, 100)$
 $root.children = []$

1) Insert (35,40)

```
algorithm insert(P:point, root:node) → bool
  if not inboundary(root.region, P) then
    return false
  end if
  if root.point is Null and |root.children| = 0 then
    root.point ← P
    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm
```

Null



Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

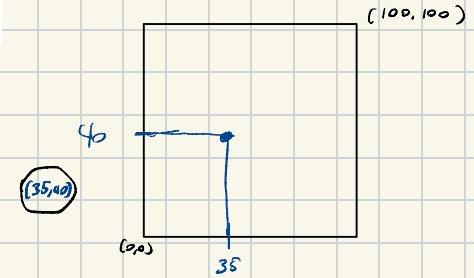
1) Insert (35,40)

```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

$root.pt = (35, 40)$

$root.region = [0, 0, 100, 100]$

$root.children = []$



Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

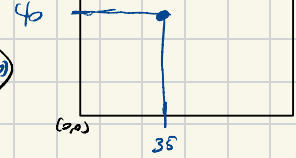
2) Insert (50,10)

root.pt = (35,40)

root.region = (0,0,100,100)

root.children = []

(35,40)



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

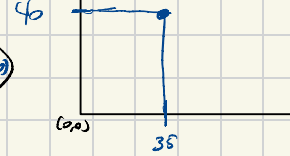
2) Insert (50,10)

root.pt = (35,40)

root.region = [0,0,100,100]

root.children = []

(35,40)



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

```
algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← (xmin + xmax) / 2
    ymid ← (ymin + ymax) / 2
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
        P ← root.point
        root.point ← Null
        for each quadrant in root.children do
            if insert(P, quadrant) then
                return
            end if
        end for
    end if
end algorithm
```

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

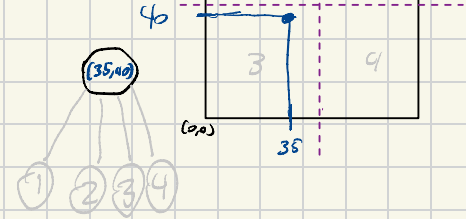
1) Insert (35,40)

2) Insert (50,10)

root.pt = (35,40)

root.region = [0,0,100,100]

root.children = []



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

```
algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← 50
    ymid ← 50
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
        P ← root.point
        root.point ← Null
        for each quadrant in root.children do
            if insert(P, quadrant) then
                return
            end if
        end for
    end if
```

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

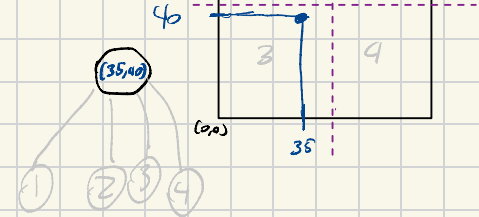
1) Insert (35,40)

2) Insert (50,10)

root.pt = (35,40)

root.region = [0,0,100,100]

root.children = []



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

```
algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← 50
    ymid ← 50
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
        P ← root.point
        root.point ← Null
        for each quadrant in root.children do
            if insert(P, quadrant) then
                return
            end if
        end for
    end if
```

What does this intuitively do?

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

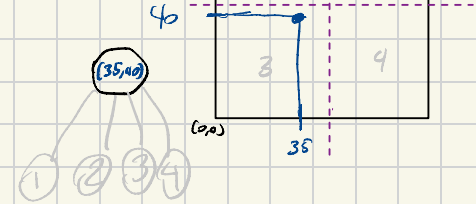
1) Insert (35,40)

2) Insert (50,10)

root.pt = (35,40)

root.region = (0,0,100,100)

root.children = []



```
algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
```

```
algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← 50
    ymid ← 50
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
```

insert root into the quadrant it belongs
(and set original root pt to null)

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

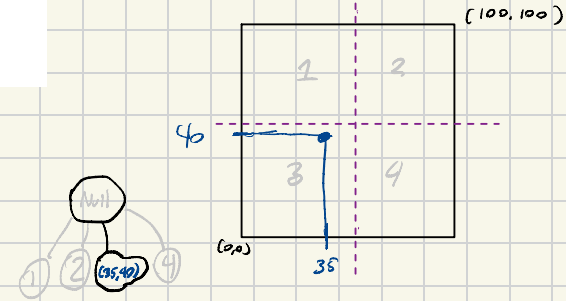
2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
    
```

```

algorithm subdivide(root:node)
    xmin, ymin, xmax, ymax ← root.region
    xmid ← 50
    ymid ← 50
    root.children ← [
        Node(Region(xmin, ymid, xmid, ymax)), //NW
        Node(Region(xmid, ymid, xmax, ymax)), //NE
        Node(Region(xmin, ymin, xmid, ymid)), //SW
        Node(Region(xmid, ymin, xmax, ymid))] //SE
    if root.point is not Null then
        insert root into the quadrant it belongs
    end if
end algorithm
    
```



• we do this on the second insert!

• still need to insert (50,10)

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

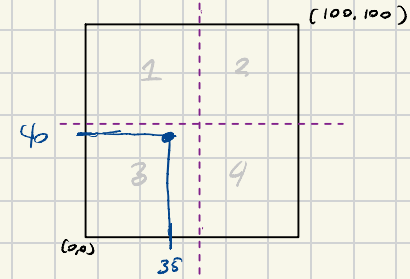
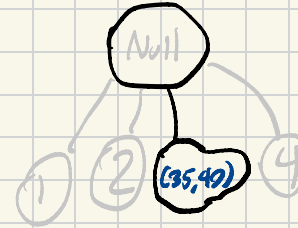
1) Insert (35,40)

2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
  if not inboundary(root.region, P) then
    return false
  end if
  if root.point is Null and |root.children| = 0 then
    root.point ← P
    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm

```



Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

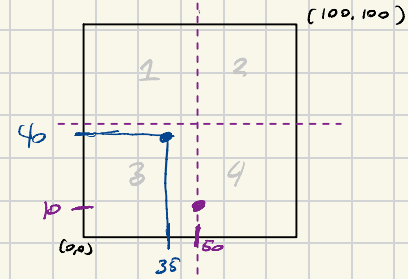
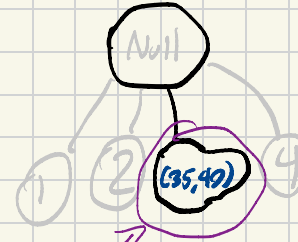
1) Insert (35,40)

2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
    
```

• insert into quadrant 3



Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

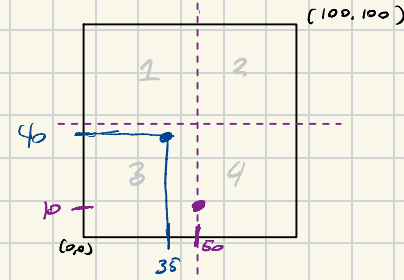
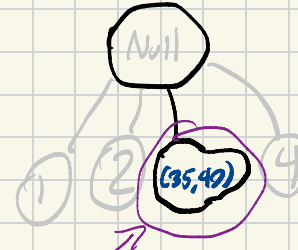
1) Insert (35,40)

2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
  if not inboundary(root.region, P) then
    return false
  end if
  if root.point is Null and |root.children| = 0 then
    root.point ← P
    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm
    
```

- insert into quadrant 3
- this node has no children, need to further subdivide, then move into new subquadrant



Question 1

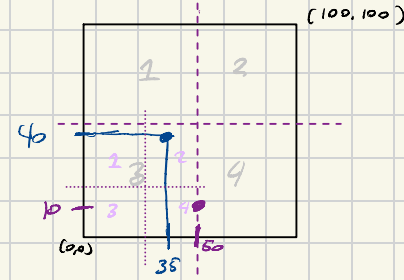
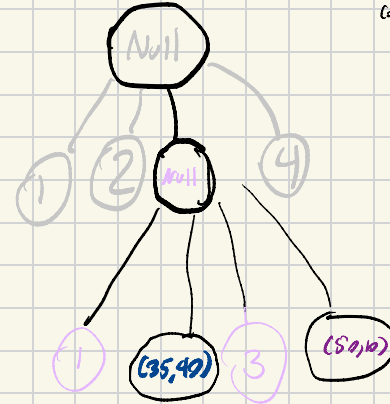
(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
    
```



- insert into quadrant 3
- this node has no children, need to further subdivide and move into new subquadrant

Question 1

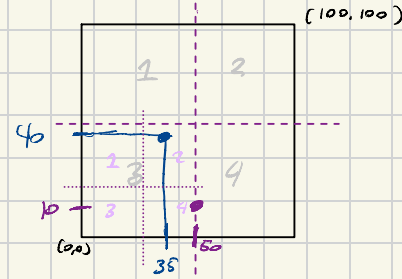
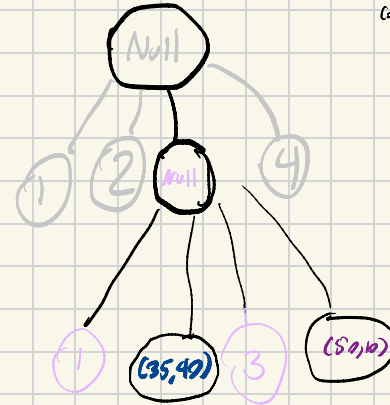
(1) (Quadtrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

1) Insert (35,40)

2) Insert (50,10)

```

algorithm insert(P:point, root:node) → bool
    if not inboundary(root.region, P) then
        return false
    end if
    if root.point is Null and |root.children| = 0 then
        root.point ← P
        return true
    end if
    if |root.children| = 0 then
        subdivide(root)
    end if
    for each quadrant in root.children do
        if insert(P, quadrant) then
            return true
        end if
    end for
end algorithm
    
```



Take away

Insert always inserts at a leaf.

Allows for $\log_4 n$ height

Question 1

(1) (Quadrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

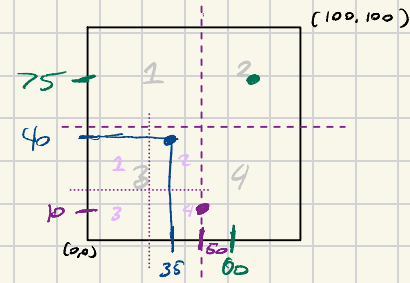
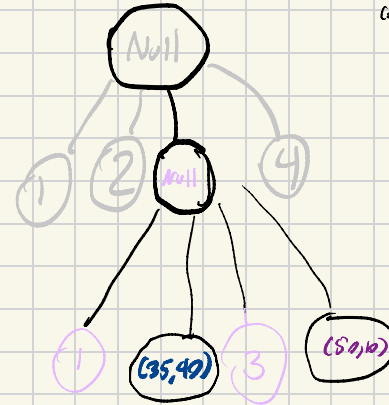
1) Insert (35,40)

2) Insert (50,10)

3) Insert (60,75)

```

algorithm insert(P:point, root:node) → bool
  if not inboundary(root.region, P) then
    return false
  end if
  if root.point is Null and |root.children| = 0 then
    root.point ← P
    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm
    
```



Question 1

(1) (Quadtrees) Insert the following points (in order) into an empty Point-region Quadtree: (35,40), (50,10), (60,75), (80,65), (85,15), (5,45), (25,35), (90,5).

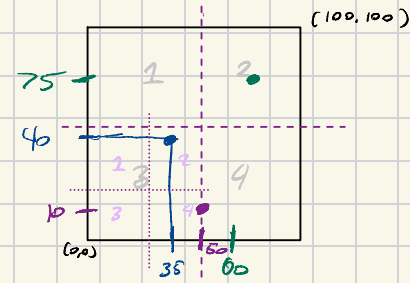
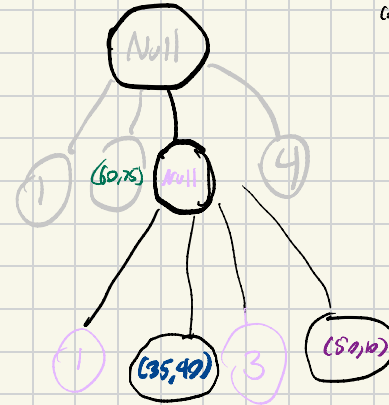
1) Insert (35,40)

2) Insert (50,10)

3) Insert (60,75)

```

algorithm insert(P:point, root:node) → bool
  if not inboundary(root.region, P) then
    return false
  end if
  if root.point is Null and |root.children| = 0 then
    root.point ← P
    return true
  end if
  if |root.children| = 0 then
    subdivide(root)
  end if
  for each quadrant in root.children do
    if insert(P, quadrant) then
      return true
    end if
  end for
end algorithm
    
```



Question 2

Consider a **QuadTree** that indexes n points uniformly distributed in a square region $[0, 1] \times [0, 1]$. The QuadTree recursively subdivides each square into four equal quadrants until each quadrant contains at most one point.

Question:

- What is the expected **depth** $D(n)$ of the QuadTree?
- What is the total number of **leaf nodes** in the tree in terms of n ?
- If we perform a **range query** for a square region of side length s , what is the expected number of leaf nodes that intersect this query region?

bst

$$2^d \approx n$$

What is the takeaway of the example for this question?

a) When we evenly sub divide into 4 regions

NOT

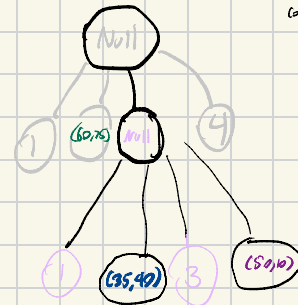
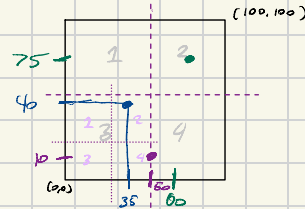
guaranteed!!

$$4^d \approx n$$

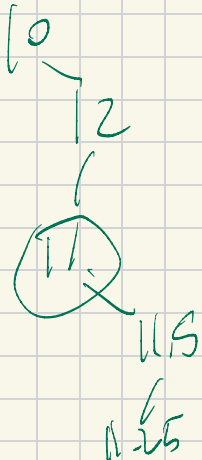
b. # non-null leaves = n

leaves $\approx n$

c. geometry with question.



$n=3$
leaves = 7



Poll

data structure fan hangout

(guest performance by fibonacci heap)

What do YOU want to see at the last practice session?

- Asymptotic analysis 50.1%
- Graph problems 40.1%

exam 1
sorting scenario

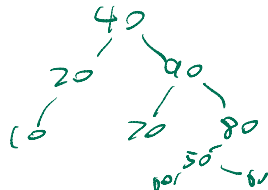
that ONE topo sort question.

b trees

hash table + Z

red black trees (deletion $\hat{=}$) pq

Do you want us to keep doing Mult. Choice or do more Free response-y questions?



2 sat reduction
dp $\hat{=}$